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Computer-Based Specifications

EDITSPEC SYSTEM MANUAL

VOLUME FOUR -- DATA HANDLER

by

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Final Technical Report

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PREFACE TO ALL EDITSPEC SYSTEM REPORTS

The purpose of the EDITSPEC system reports is to provide complete documentation to all personnel that must be involved with the EDITSPEC system. Such personnel include managers, specification writers, typists, computer systems analysts, and computer programmers. Each personnel group requires different documentation. The reports required and the order of reading are shown in the table below.

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ABSTRACT

This report provides computer programmers with documentation for a machine-independent scientific and engineering data-handler. The report discusses design concepts, subroutine functions, application commands, error messages, program conversion, test routines, and procedures for changing the data-handler. The data-handler has been designed to minimize the resources required for its conversion and operation.

FOREWORD

This investigation was performed for the Directorate of Military Construction, Office of the Chief of Engineers (OCE), under Project 4A762731AT41, "Design, Construction, and Operation and Maintenance Technology for Military Facilities"; Task Tl, Development Work Unit 009, "Computer-Based Specifications." The applicable QCR is 1.10.001. The OCE Technical Monitor was William Darnell.

The basic computer programs were written by Multi-Systems Inc., Cambridge, Mass., under contract DACA 23-75-C-0003, and modified by Hans Wegener and Jayant Krishnaswamy of CERL.

The study was performed by the Management Systems Branch (Dr. O. E. Rood, Jr., Chief), Facility Acquisition and Construction Division (Mr. E. A. Lotz, Chief), U.S. Army Construction Engineering Research Laboratory (CERL).

COL Louis J. Circeo is Commander and Director of CERL, and Dr. L. R. Shaffer is is Technical Director.

LIST OF REVISIONS TO THIS DOCUMENT

Revision Number

Date <u>Revised</u>

Reason for Revisions

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1 Design Concepts Introduction Definitions Data Storage Modes Data Management Imp	lementation	11122
2 Applying the Data Hand	dler 3	33
3 Application Command Son Definitions Learning the Command Command Descriptions	ds 😜	4
4 Subroutine Functions	6	2
5 Error Messages	6	3
6 Program Conversion	6	5
7 Documentation Tapes	6	7
8 Test Routines	,	9
9 Changes to Data-Handle	er Programs 7	ď
APPENDIX A Data Handler APPENDIX B Subroutine Fo	•	

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EDITSPEC SYSTEMS MANUAL VOLUME IV: DATA HANDLER

1 DESIGN CONCEPTS

Introduction

This data-handler is a feasible method by which a programmer can create data structures on disk without knowing record sizes, data locations, etc. This direct access system provides in-core data handling so that physical disk accesses are minimized.

The programs consist of several data storage and retrieval functions and a variety of other utilities which create and delete files, inform the user of file status, and establish or remove access restrictions. The programs are written mostly in ANSI FORTRAN.

Definitions

DATA SET: a named, independently accessible collection of data on a direct access storage device, created and maintained by the host operating system; called a file by some hardware manufacturers.

PHYSICAL RECORD: a physical portion of a data set usually of fixed length, which is accessible by means of one input/output operation the computing system.

FILE: a named, logical subdivision of a data set.

LOGICAL RECORD: a logical portion of a file of variable length, which is accessible through one application of the data-handler; usually referred to simply as "record."

In essence, the data-handler maintains a logical structure of named files composed of variable-length, randomly accessible records by manipulating physical records in one or more data sets.

APPLICATION PROGRAM: The set of subroutine that are using the DH for data base management.

Data Storage Modes

Two storage modes control how data is stored in the physical records on disk: CHAIN and PACK.

In CHAIN mode, the data-handll store with each logical record the record identifiers of the preceding and following logical records in the file. The logical records must therefore remain in the order in which they are written. Chain pointers are available to the programmer who is reading or updating logical records. The pointers are always updated whenever a logical record is added or lengthened.

In PACK mode, the data-handler routines will fill physical records to capacity and, when necessary, segment logical records which span two or more physical records.

CHAIN and PACK modes may be specified independently so that four combinations of options are available:

- 0 neither CHAIN nor PACK
- 1 CHAIN only
- 2 PACK only
- 3 CHAIN and PACK

The DH recognizes only the first definition of this option. Any redefinitions are ignored by the DH.

Data Management Implementation

The "data handler" (DH) is a general-purpose, direct-access data management system. The DH appears to an application program as a set of subroutines. The DH is written largely in ANS FORTRAN. The DH is designed as an interface between an application program and FORTRAN direct-access input/output (I/O) facilities (Figure 1). This provides a two-fold advantage. Direct-access I/O can be performed (1) at a logical level and (2) in a machine-independent manner.

The DH deals with only direct-access data sets (Figure 2). The contents of a dataset must be understood from three different view points:

- (1) host system
- (2) data handler
- (3) application program
- a. <u>Host System</u>. Data sets consist of physical records.
- b. <u>Data Handler</u>. Data sets consist of files. Files consist of physical records. Physical records may also be referred to as

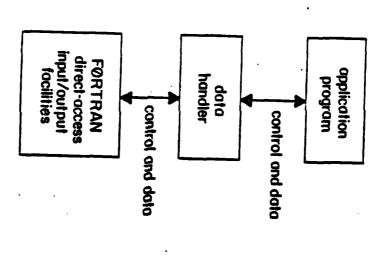
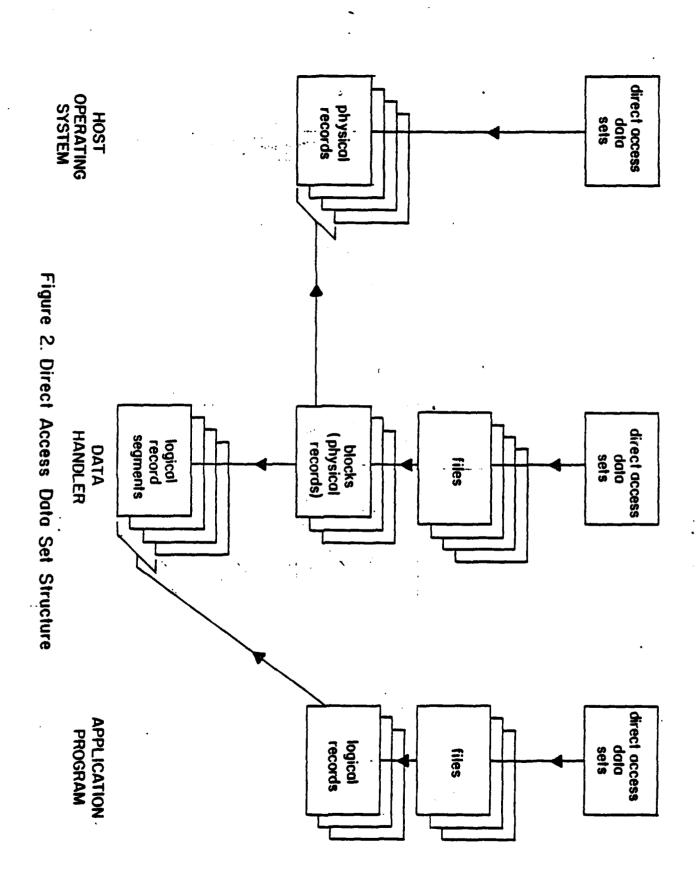


Figure I. Data Handler As Interface



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blocks. Blocks (physical records) are composed of logical record segments.

c. <u>Application Program</u>. Data sets consist of files. Files consist of logical records.

Several cross-relationships should be noted:

- (1) The logical record of the application program may actually be stored within one or more data handler logical record segments.
- (2) A data handler block may be written as one or more operating system physical records.

Program Variable Naming Conventions

The names used for variables within the data handler follow the following logical order:

- (1) All variable names are five letters in length.
- (2) The first letter options are:
 - a. K for contents of
 - b. I for index to
 - N for number of
- (3) The second and third letters options are:
 - a. PR for physical record
 - b. LR for logical record
 - c. DS for data set
 - d. OF for open file
 - e. BT for block table
- (4) The fourth letter options are:
 - a. H for header
- (5) The fifth letter options are:
 - a. S for status
 - b. P for previous

- c. F for following
- d. N for number

Data Set Structures

The first two blocks of a direct-access data set maintained by the DH contain three independent, special-purpose LR segments (Figure 3). These blocks do not belong to any file, and the segments are known as the data set header (DSHED), the file directory (FD), and the free space table (FST). The DSHED and FD are contained in the first block, while the FST is contained in the second.

The DSHED contains various information about the data set and host computer system (Figure 4). See the discussion of data set table entry below for the meanings of KDSSU, KDSPR, KDSBI, KDSFS, and KDSFD. See Table 1 for the meanings of NBSGN, NBBLK, DHFLG, ØPSYS(2), and DHVER. The time is given in the form hh:mm:ss, and the date is given left-justified and blank-padded in the form dd-mmm-yy; they are the time and date when the data set was initialized by the DH. The rest of the elements in the DSHED are the same as identically-named elements in blank COMMØN.

The FD contains an entry for each file that is contained in the data set. An FD entry is identical to an open file table entry. FD entries are numbered consecutively beginning with one.

The FST is a bit map; each bit indicates whether (on) or not (off) the block corresponding to that bit has been allocated to a file in the data set. Both the FD and FST are extended when necessary to additional LR segments.

The remaining blocks are used to store data. The data is organized into dynamic collections of blocks known as files. One block can belong to only one file.

Block Structures

Blocks within a dataset are sequentially numbered starting with one. A block can belong to only one file. The block number is the sequential number assigned to a block. A block is composed of three sections.

The first three SU's in each block are known as the physical record (PR) header (figures). The first SU contains the contents of the physical record header status (KPRHS). The high-order (left-hand) half of KPRHS is the number of logical record segments in the block, while the low-order half is the number of unused SU's at the end of the block. The second SU KPRHP is the block number of the block that was previously written for the file to which this block belongs; it is zero if this block is the first one in the file. The

Data Set Header File Directory	Free Space Table		•
ck l	Block 2		
Block	Bloc		-

Figure 3. Data Set Layout

				_P			γ -		. ,				,		
KDSSU	KDSPR	DHFLG	NBSGN	NBBLK	OPSYS(2)		DHVER	NBW	NCW	NAW	NBC	NCO	NBU	NAU	
						<u> </u>									

dote (3SU's) | dd-mm-yy

Initialized

KDSFS

zero

KDSBI

KOSFD

time (2 SU's) hhimmiss

Initialized

LFIWU

LFIW

NWPSU

RVRSE

NBU2

NC2U

Figure 4. Data Set Header

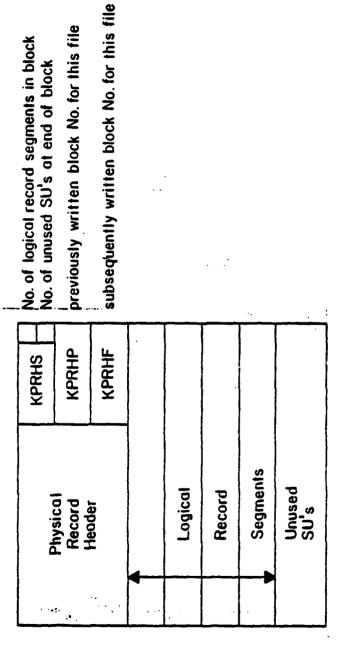


Figure 5. Block Layout

third SU KPRHF is the block number of the block that was subsequently written or is next to be added to this file. Hence, the blocks that constitute a file are chained together in both directions.

The application programs logical record may be divided by the DH into segments for storage. These segments are known as logical record segments. The block will contain one or more logical record (LR) segments. There are three reasons why an LR may be represented by more than one segment: (1) the LR is too long to fit in one block; (2) the LR is too long to fit into the unused portion of any block belonging to the file for which this LR is written, and the pack option is in effect for this file; (3) the LR cannot be contiguously extended due to lack of space in the block.

The remaining portion of the block contains unused SU's.

Logical Record Segment Structure

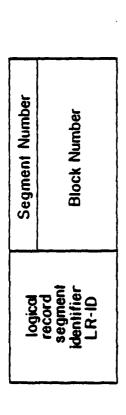
Associated with each LR segment is an SU quantity known as the LR segment identifier (ID). The high-order quarter of this quantity is the segment number, which is the position of the segment within the block, relative to the beginning of the block; segments are numbered beginning with zero. The low-order three-quarters of the LR segment ID constitute the block number. If the LR segment is the first or only segment of the LR, then the LR segment ID is often called simply the LR ID.

The first several SU's in each LR segment are known as the LR segment header (Figure 6). The extreme high-order portion of KLRHS contains three status bits known as the continue, segment, and delete bits, whose positions are given by the COMMON variables NBCON, NBSEG, and NBDEL, respectively (Table 1). Bits within an SU are numbered from right to left beginning with zero. The continue bit indicates whether (on) or not (off) this LR segment is continued to a subsequent segment. The segment bit indicates whether (on) or not (off) this segment is continued from a previous segment. The delete bit indicates whether or not this LR has been deleted. The four loworder bits in the high-order half constitute the number of free characters in the last SU. The low-order half of KLRHS is the number of SU's in the LR (segment) including the header. The segments that represent an LR are chained together in both directions, but the LR's that constitute a file are chained only if the chain option is in effect for the file (Figure 7).

KLRHP contains the following:

No Chain Option

(1) If this LRS is the first one in the file, zero.



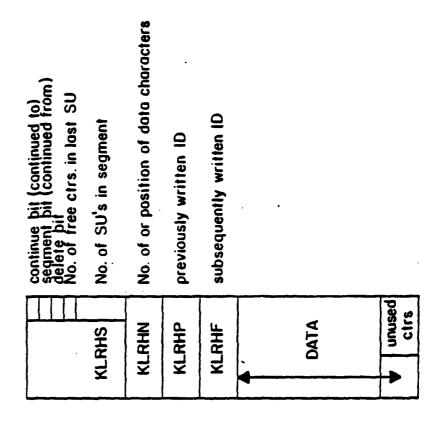


Figure 6. Logical Record Segment Layout

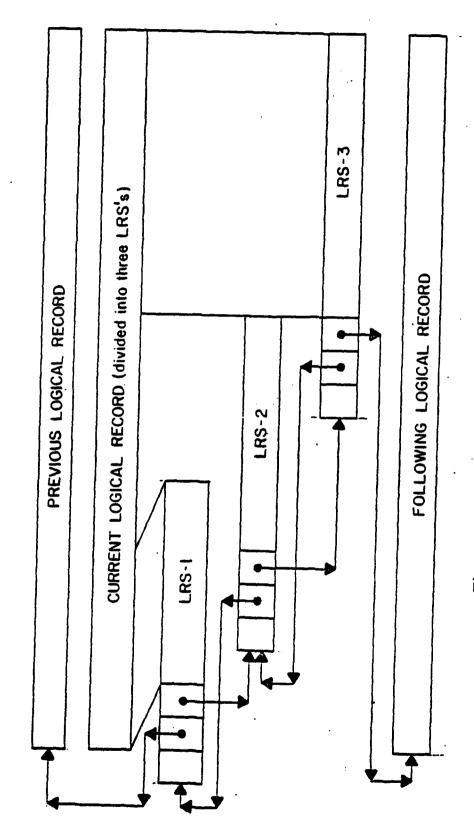


Figure 7. Logical Record Segment Chain

TABLE 1
FIXED AREA

Variable/ Array	Position in COMMCN	Description Value
		1. SYSTEM WIDE
	1	
	2	
	3	•
	4	
		*** Machine Characteristics ***
NBW	5	Number of bits per word
NCW	6	Number characters in a word
NAW	7	Number of addresses in a word
NBC	8	Number of bits per character
NCU	9	Number of characters in a : "standard unit" (approximately 32 bits long)
NBU	10	Number of bits per standard unit
NAU	11	Number of addresses in a standard unit
NBU2	12	One-half of NBU
NC2U	13	Number of characters in a double precision variable (about 8-10 chars)
RVRSE	14	<pre>Flag telling order of characters: 0 = left to right 1 = right to left</pre>
NWPSU	15	Number of real words per standard unit: 0 = two SU's per real (e.g. IBM 1130) 1 = one SU per real (e.g. IBM 360) 2 = two real's per SU

TABLE 1 (CONT'D)

Variable/ Array	Position in COMMON	Description	Value
LFIW	16	Number of bits in a FORTRAN integer	
LFIWU	17	Number of useable bits in a FORTRAN integer	
	18		
	19		
	20		
	21	•	
. •	22		
•	23	·	
		*** Standard Devices ***	
ICRD	24	Card reader logical unit number	
IPRT	25	Line printer logical unit number	•
IPCH	26	Card punch logical unit number	
ITAP	27	Magnetic tape logical unit number	
ITRMI	28	Terminal input logical unit number	•
ITRMO	29	Terminal output logical unit number	
JTRMI	30	Second terminal input logical unit number	
JTRMO	31	Second terminal output logical unit number	
IPLT	32	Plotter logical unit number	
IGRAI	33	Interactive graphics input logical unit	
IGRAC	34	Interactive graphics output logical unit	

TABLE 1 (CONT'D)

Variable/ Array	Position in COMMON	Description	Value
	35		
	36		
	37		
		*** Constants and Masks ***	
TRU	. 38	Value of logical true (largest positive integer)	
FALS	39	Value of logical false (largest negative integer)	
ZEROS	40	A word of all "zeros"	
ONES	41	A word of all "ones"	
FORM1	42	Token format mask (1,2,1)	
FORM2	43	CDL dictionary format mask (2,2)	
FORM3	44	CDB pointer format mask (3,1)	
	45		
	46		
	47	•	
	48		
	49		
		*** Executive Data Area ***	
SCA	50	Subscript of start of communications area (=176)	
ATLID	51	LRID of active task list	
SCRFL	52	Name of executive work file	
SCRFL	53	Name of executive work file	

TABLE 1 (CONT'D)

Variable/ Array	Position in COMMON	Description	Value ————
	54		
		*** Lexical Analyzer Data Area ***	•
NXCHR	55	Next character in input buffer	
IDFT	56	Data field type	
ISTD	57	Start of data field	
LSTR	58	Length of string or an integer value	
DCOM	59	Command (up to 8 characters) or a	
DCCM	60	double precision value	
MLCA	61	Maximum length of user communications area (for referencing in a CDB)	
SYDID	62	LRID of system command dictionary	
DICID	63	LRID of current command dictionary	
CDBNM	64	Name of current CDB	
CDBNM	65	Name of current CDB	
CDBID	66	LRID of current CDB	
SYMID	67	LRID of current symbol table	
SVTID	68	LRID of current standard value table	
NXELM	. 69	Pointer to next CDB element	
RELOC	70	Subsystem common relocation factor (must be a positive number)	
CSTID	71	LRID of modifier condition stack	
ATID	72	LRID of ID wait list	
	73		•
	74		

TABLE 1 (CONT'D)

Variable/ Array	Position in COMMON	Description	Value
	75		
	2. DATA	HANDLER INTERNAL VARIABLE AND ARRAYS.	
IPOCL	76	Position in COMMON of beginning of pool area	201
IPLEN	77	Length of pool area	9800
NBLMX	78	Maximum length of a block	256
NBSGN	79	Number of bits in a segment number	8
NBBLK	80	Number of bits in a block number	24
ILLRH	81	Length of an LR (segment) header	4
ILRHS	82	 Position in LR (segment) header of KLRHS 	1
ILRHN	83	Position in LR (segment) header of KLRHN	2
ILRH	84	Position in LR (segment) header of KLRHP	3
ILRHF	85	Position in LR (segment header of KLRHF	4
ILPRH	86	Length of a PR header	3
IPR4S	87	Position in PR header of KPRHS	. 1
IPRHP	88	Position in PR header of KPRHP	. 2
IPRHF	89	Position in PR header of KPRHF	3
IBTST	90	O'th displacement of BT	(compute

TABLE 1 (CONT'D)

Variable/ Array	Position in COMMON	Description	Value
NEBT	91	Number of BT entries	(computed)
IELBT	92	Length of a BT entry	273
ILBTH	93	Length of a BT header	17
IBTNS	94	Position in BT header of KBTNS	1
IBTBL	95	Position in BT header of KBTBL	2
IBTQ	96	Position in BT header of KBTQ	3
IBTPP	97	Position in BT header of KBTPP	4
IBTL	98	Position in BT header of KBTL	5
IBTTR	99	Position in BT header of KBTTR	6
IBTCN	100	Position in BT header of KBTCN	7
IBTDP	101	Position in BT header of KBTDP	8
IBTNC	102	Position in BT header of KBTNC	9
IBTSW	103	Position in BT header of KBTSW	10
IBTFN	104	Position in BT header of KBJFN	11
IBTFI	105	Position in BT header of KBTFI	15
IBTPI	106	Position in BT header of KBTPI	16

TABLE 1 (CONT'D)

Variable/ Array	Position in COMMON	Description	Value
IBTLN	107	Position in BT header of KBTLN	17
IBTLG	108	O'th displacement of BT entry referenced as result of most recent call of DKGET	(variable*)
IBTPQ	109	O'th displacement of BT priority queue	(computed)
IBTOP	110	Position in BT of entry with highest priority	(variable**)
IBTBT	111	Position in BT of entry with lowest priority)	(variable***)
IOFST	112	O'th displacement of OFT	(computed)
NEOFT	113	Number of OFT entries	125
IELOF	114	Length of an OFT entry	9
IOFFN	115	Position in OFT entry of KOFFN	5
IOFFL	116	Position in OFT entry of KOFFL	6
IOFLL	117	Position of OFT entry of KOFNB	7
IOFNE	118 .	Position in OFT entry of KOFNB	8
ICFLB	119	Position in OFT entry of KOFLB	8
ICFFD	120	Position in OFT entry of KOFFD	9
IDSST	121	O'th displacement of DST	200
IELDS	122	Length of a DST entry	12

TABLE 1 (CONT'D)

Variable/ Array	Position in COMMCN	Description	Value
IDSLU	123	Position in DST entry of	1
IDSFN	124	Position in DST entry of KDSFN	2
IDSPR	125	Position in DST entry of KDSPR	4
IDSPB	126	Position in DST entry of KDSPB	5
IDSPL	127	Position in DST entry of DKSPL	6
IDSFD	128	Position in DST entry of KDSFD	. 8
IDSFS	129	Position in DST entry of KDSFS	9 ·
IDSSU	130	Position in DST entry of KDSSU	7
ILDSN	131	Length of a data set name	2
OPSYS(2)	132	Name of host operating system	'IBM OS'
DHVER	134	DH version	'V1F '
DHFLG	135	OH flag	י אם' י
NBCON	136	Position in KLRHS of continue bit	29
NESEG	137	Position in KLRHS of segment bit	30
NBDEL	138	Position in KLRHS of delete bit	31
NBWIN	139	Position in KBTSW of written-into bit	0

TABLE 1 (CONT'D)

Variable/ Array	Position in COMMCN	Description	Value		
NBRIT	140	Position in KBTSW of write-in-progress bit	1		
NBRED	141	Position in KBTSW of read-in-progress bit	2		
NBOPN	142	Position in KDSLU of open bit	16		
NBCLD	143	Position in KDSLU of old bit	17		
NBSAV	144	Position in KDSLU of save bit	18		
NBSHR	145	Position in KDLSU of share bit	19		
NBCHN	146	Position in KOFFD of chain bit	C		
NBPAK	147	Position in KOFFD of pack bit	1		
NBPRO	148	Position in KOFFD of protect bit	4		
NBNDS	149	Bit position in KCFFD of beginning of data set number	8		
КО	150	Zero	C		
K1	151	One	1		
K2	152	Тwo	2		
IDSBI	153	Position in DST entry of KDSBI	10		
IDSWN	154	Position in DST entry of KDSWN	11		
IDSVV	155	Position in DST entry of KDSWD	12		

TABLE 1 (CONT'D)

Variable/ Array	Position in COMMON	Description	Value
HEDLR(4)	156	Work area for an LR (segment) header	(variable)
DST(12)	156	Work area for a DST entry	(variable)
DATA(17)	160	Work area for a BT header	(variable)
	3.	DATA HANDLER COMMUNICATIONS.	
STASK	176	Sending task name (RSX-11D scratch area)	
LCA	177	Length of user communications area transmitted by last/next message set	
TRACE	178	Executive trace switch: each bit, when on, indicates a type of tracing to be printed on unit IPRT. Bit (Value) Item Printed O 1 Digitize expansions 1 2 Tokens created 2 4 Macro expansions 3 8 Macro argument	•

TABLE 1 (CCNT'D)

Variable/ Position Array in COMMON		Description					
EXST1	179 .	Executive status switch 1: Each bit, when on. signifies the presence of some exceptional condition. Bit (Value) Meaning O 1 Return to executive 1 2 ID labels required 2 4 ID wait list processing 3 8 Inhibit user program execution (system control) 4 16 Inhibit user program					
÷.;		execution (user control) Next record continuation of preceding command Repeat tabular processing End of file on input Input error occurred (CI) User program signaled error Louis 4096 Interactive mode Interactive mode Interactive mode					
	180						
PSTID	181	LRID of program/CDB stack					
TORID	•182	LRID of token list					
TVLID	183	LRID of token value table					
NXTKN	184	Next token to be processed					
	185						
INCMD	186	Logical unit number from which to read commands *** Data Handler Arguments ***					
MAXGP	187	Maximum number of characters per GET/PUT					

TABLE 1 (CONT'D)

Variable/ Array	Position in CCMMON	Description	Value
NDS	188	Data set number	
FILNM	189	File name (up to 8 characters)	
FILNM	190	File name (up to 8 characters)	
LULID	191	LRID of next (or first) record in file	
PREID	192	LRID of preceeding (or last) record	
LRLEN	193	Logical record length in characters	
FSTAT	194	File status (from DKFIL)	
.NDA	195	Maximum number of direct access data set	
IDHTR	196	Data handler trace switch	
IDHER	197	User error handling switch	
IDKER	198	Data handler error number	
SSYSN	199	Subsystem name (up to 8 characters)	
SSYSN	200	Subsystem name (up to 8 characters)	

^{*} initialized to value of IBTST
** initialized to 1
*** initialized to value of NEBT

(2) If this LRS is not the first LRS, LRS ID of the LRS written previously for this LR.

Chain Option

- (1) If this LRS is the first one in the file, zero.
- (2) If this is the first LRS of a LR, LR ID for the first LRS of the previously written.

KLRHF contains the following:

No Chain Option

- (1) If LRS not the last LRS in LR, LRS ID of the subsequent LRS.
- (2) If LRS is the last LRS in LR, or the point option is not in effect, zero.

Chain Option

(1) If LRS is the last LRS in CR, LR ID for the first LRS of the subsequent record.

The actual data is stored next followed by the unused characters in the last SU.

Program Common Structure

The great majority of internal variables and arrays used by the DH are located in blank CØMMØN (Figure 8). Common consists of a fixed area and a pool (dynamic) area. The first sections in the fixed area contain system-wide and DH communications variables and arrays. The second section contains the DH internal variables and arrays. All three sections are documented in Table 1. All positions and lengths in Table 1 are given in terms of SU's. All elements in this section that need to be initialized are done so during DH initialization.

The pool area contains the four DH tables: the data set table (DST), the open file table (ØFT), the buffer table (BT), and the BT priority queue. There is an unused area at the end of common.

The DST contains an entry for each data set that is currently in use by the DH (Figure 9). The high-order half of KDSLU contains four status bits known as the open, old, save, and share bits, whose positions are given by the COMMON variables NBOPN, NBOLD, NBSA, and NBSHR, respectively (Table 1). The open bit indicates whether or not the data set is open to the host operating system. The old bit indicates whether or not the data set existed before the current

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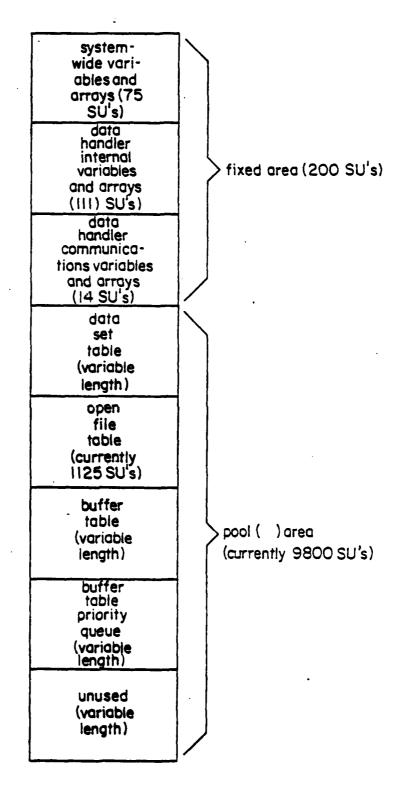


Figure 8. Common Layout

Open status bit Old status bit Save status bit	Shore status bit Logical unit No.	Data set name		Phsical record size in SU's	No. whole PRs per block	No. SU's last partial PR in block	block size in SU's	entry No. next free FD entry	No. of blocks used in data set	No, of blocks initialized in dato so	position of KDSWN in FST	
KOSLU		KDSFN(2)		KDSPR	KDSPB	KOSPL	KDSSU	KDSFD	KDSFS	KDSBI	KDSWN	KDSWD
			Set in Use									
		i	Set in Use									

Figure 9. Data Set Table Layout

execution of the DH. The save bit indicates whether or not the data set is to be saved at the end of the current execution. The share bit indicates whether the data set can be shared with another, concurrently-executing DH task. The low-order half of KDSLU is the logical unit number.

KDSFN is the data set name, KDSPR is the PR size (in SU's), KDSPB is the number of whole PR's per block, KDSPL is the number of SU's in the last (partial) PR of a block if the block does not end on a PR boundary (it is zero otherwise), KDSSU is the block size (in SU's), KDSFD is the entry number of the next free FD entry, KDSFS is the number of blocks initialized in the data set, and KDSWN is the postion of KDSWN within the FST.

The ØFT contains an entry for each file that is currently in use for output by the DH (Figure 10). KØFFN(4) is the file name. KØFFL and KØFLL are the LR ID's of the first and last LR's, respectively, that were written for the file. KØFNB is the block number of the next block to be added to the file. KØFLB is the block number of the last block that was written for the file. The high-order half of KØFFD is the entry number of the FD entry associated with this file. The bits in the low-order half of KØFFD constitute the data set number (DST entry number) of the data set to which the file belongs. The extreme low-order portion of the low-order half of KØFFD contains three status bits known as the chain, pack, and protect bits, whose positions are given by the CØMMØN variables NBCHN, NBPAK, and NBPRØ. The chain and pack bits indicate whether or not the chain and pack options are in effect for the file. The protect bit indicates whether or not write protection is in effect for the file.

The BT contains an entry for each block of a data set that is currently in main storage (Figure 11). KBTNS is the data set number (DST entry number) of the data set to which this block belongs. KBTBL is the block number of the block. KBTQ is the segment number of the segment which was requested when the block was read into main storage, and KBTPP is the position in CØMMØN of the SU immediately proceeding the beginning of the segment. KBTL and KBTTR are the segment and block numbers, respectively, of the LR that was most recently referenced in this BT entry. KBTCN is the position within this LR of the first (data) character (of a set) that was requested. KBTDP is the position in CQMMQN of the SU that contains this character. KBTNC is the number of characters in the first segment of this LR. KBTSW contains three status bits known as the written-into, write-in-progress, and read-in-progress bits, whose positions are given by the COMMON variables NBWIN, NBWIN, NBRIT, and NBRED, respectively. The written-into bit indicates whether or not the block has been modified in this BT entry. The write-in progress bit indicates whether or not an output operation is in progress for the block. The read-in-progreas bit indicates whether or not an input operation is in progress for the block. KBTFN(4) is the name of the file to which the block belongs. KBTFI and KBTPI have the same values as KLRHF and

File Name		•	LRS-1 First LR	LRS-1 Lost LR	block No. next block to add	block No. last block added	entry No. in FO data set number chain pack protect
KØFFN(4)			KOFFL	KØFLL	KØFNB	KØFLB	KØFFD
	File Two	•					
	File One						

Figure 10. Open File Table/File Directory Entry Layout

٢				٦\							
			KBTNS	`	data	set No	oDS	T			
	-		KBTBL		block	ι No.					
	Buffer No. I	Buffer No. 2	KBTQ		segn	nent No	o. req	juested			
			KBTPP		su po	noitiec	proce	eeding s	egmer	nt	
			KBTL		most	recent	lly ref	ferenced	d segn	nent	No.
			KBTTR		n	It		le .	11		block No.
			KBTCN		first	data c	harac	cter posi	ition		
-		!	KBTDP		}			st data (
			KBTNC		head	er	•	st segme	ent		
			KBTSW		write	en into in prog	gress	bit			
			KBTFN(4)		file n	in prod ame	gress	DIT			•
			·								
		Ī	KBTFI		KLRH	dF for	KBT	L			
			KBTPI		KLRH	IP for	KBTI	ΓR			
			KBTLN	1	No. ct	trs. in	LR				
				>							
			data		buffe	r (curr	ently	256 St	J's)]		
				1							

Figure 11. Buffer Table Layout \mathcal{SC}

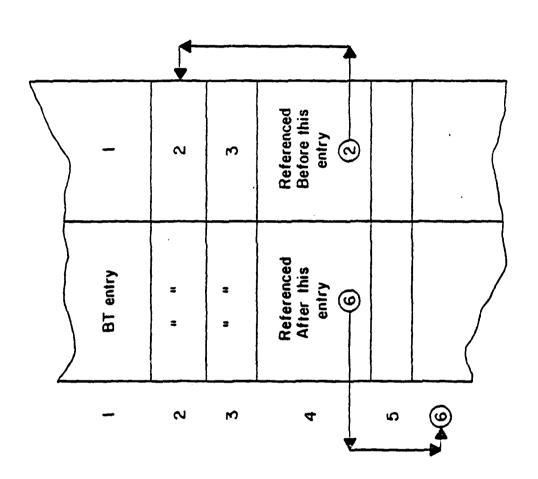


Figure 12. BT Priority Que Layout

KLRHP, respectively, in the LR header of the LR given by KBTL and KBTTR. KBTLN is the number of characters in the LR.

The BT priority queue contains an SU entry for each BT entry (Figure 12). The high-order half of the i'th BT priority queue entry is the entry number of the BT entry that was referenced immediately after the i'th BT entry. The low-order half of the i'th BT priority queue entry is the entry number of the BT entry that was referenced immediately before the i'th BT entry. Hence, the BT priority queue imposes a most-recently-referenced-to-least-recently-referenced ordering on the BT entries. A more recently referenced BT entry is said to have a higher priority than a less recently referenced BT entry. A block in a higher-priority BT entry will be retained in main storage longer than a block in a lower-priority one.

The DH includes a routine which produces a dump of blank CØMMØN. All items are identified in the same manner as they are in this documentaion. The values are given in decimal digits and characters as applicable except for KDSLU, KØFFD, KBTSW, the buffer contents of the BT entries, and the BT priority queue entries, which are given in hexadecimal digits. An annotated dump is given in Appendix A.

Definitions

- Data Set Block Fixed-sized unit of information stored on a data set. The actual information read (or written) from the disk and to the local program I/Ø buffer area.
- Program I/Ø Buffer Fixed-sized unit of storage in the local program space to be used to store data read from (to be written to) disk storage.
- Physical Becord Fixed-sized unit of information measurement which must be equal to or less than the size of the data set block.

2 APPLYING THE DATA HANDLER

The use of the data handler (DH) within a computer system application program requires the review and application of two data handler routines. The first routine, DKNIT, initializes the data handler, allows the system programmer to define the system parameters to the data handler, and builds the necessary tables for DH execution.

The system programmer must review the definition of each parameter and variable in this routine and define the values with respect to the actual operating system characteristics.

This routine must be the first routine called by the application system.

4:4:

in:

When all processing has been completed the application system should call DKXIT to logically close all DH managed data sets.

During the period after the call to DKNIT and before DKXIT all application programs can utilize the DH application routines to read and/or work information from/to the DH datasets.

3 APPLICATION COMMAND SUBROUTINE DESCRIPTION

<u>Definitions</u>

Several parameters are kept in COMMON, rather than being passed as arguments. These include:

FILENAME - the name of the file being operated on; it contains up to 16 characters.

NDS - the data set number, usually 1,2,3, or 4.

PREID - the logical record identifier of the preceding record in the file.

FOLID - the logical record identifier of the following record in the file.

LRLEN - the logical record length, in characters.

FSTAT - the file status (see DKFIL 4.8).

Learning the Commands

The logical order for learning the commands and functions is to read the descriptions in the following order:

1. Access to Data.

- 1. LOK/UNLOK to obtain controlled access to information.
 - 2. Data Sets.
- 1. DKNNDS To find a data set number.
- 2. DKDSN To find data set name
- 3. DKSET To obtain general information.
 - 3. Files.
- DKOPN To create and/or open a file for writing.
- 2. DKPUT To write information.
- 3. DKINS To write (insert) information
- 4. DKGET To read information
- 5. DKDEL To delete information.
- 6. DKRNM To rename the file
- 7. DKFIL To obtain general information
- 8. DKLEN To obtain the length of a record.
- 9. DKLOS To disallow writing

To delete a specific record within a file or to delete a complete file.

GENERAL FORM

CALL DKDEL (RECID)

where RECID

FIELD

is a zero or the record id for an existing record. (INPUT)

DEFAULT

FIELD OPTIONS

1.	RECID	 valid record id zero (0) to delete all records in the file FILENAME	No Default

OPTIONS

SPECIAL NOTES

The programmer must set variables NDS and FILENAME in common before the subroutine is called.

EXECUTION PROCEDURE

The system will delete the record specified from the data set NDS. If the RECID is zero, the entire file defined in FILENAME will be deleted from data set NDS.

To find the dataset name when the dataset number is known.

GENERAL FORM

CALL DKDSN(DSNM)

where

is returned as the dataset name (output)

FIELD OPTIONS - None

SPECIFIC NOTES

The programmer must set variables NDS and FILENAME in common before the subroutine is called.

EXECUTION PROCEDURES

The system will locate the data set specified by NDS and return the dataset name. If the data set is not found, the routine will return the next data set with a larger NDS valve. NDS will be changed to point to the new data set. If no more data sets are available, NDS and DSNM will be zeroed.

To obtain the general information about a specified file.

GENERAL FORM

CALL DKFIL

FIELD OPTIONS - None

SPECIAL NOTES

The programmer must define NDS and FILENAME in common before this sub-routine is called.

EXECUTION PROCEDURES

FSTAT is returned in common as

- 1 if the file does not exist
- 2 if the file is closed
- 3 if the file is open

PREID is returned in common as the record id of the first record in the file.

FOCID is returned in common as the record id of the last record in the file.

To obtain information from a file on disk storage.

GENERAL FORM

CALL DKGET (RECID, AREA, L1, L2, NC)

where	RECID	is the logical record identifier. (INPUT)
	AREA	is the location into which the information is to
	L1	be placed. (OUTPUT) is the first character to be read
	L2	from the record. (INPUT) is the last charac ter to be read
	NC	from the record. (INPUT) is the actual number of characters read from
		the record. (OUTPUT)

FIELD OPTIONS

FIELD	OPTIONS	DEFAULT
1. RECID	a. valid record	No Default

FIELD OPTIONS

FIE	ELD	OPTIONS	DEFAULT
2.	AREA	a. one computer word, variable name b. the first computer word to be used to hold the data. An element of a dimensioned array.	No Default
3.	L1	 a. less than one or greater than the length of the logic record, it is interpreted as being equ to one. b. the first character to be read. 	al . al
4.	L2	a. greater than the number of character in the specified logical record or i less than Ll, it is interpreted as being equal to the last character in the record. b. the last character to be read. c. zero - is interpreted as being equal to the last character in the record.	s f g
5.	NC .	Actual number of of characters read.	No Default

SPECIAL NOTES

The programmer must set NDS in common before a call to this subroutine. The programmer can read an entire logical record of unknown length by setting L1 and L2 equal to zero. The programmer must define RECID, AREA, L1 and L2.

EXECUTION PROCEDURES

The system will read the requested information (L1, L2) from the requested record RECID located on data set NDS. The data will be stored starting with word AREA. The system will return NC as the number of characters actually placed into AREA.

To insert a record at the beginning of a point file or to insert a record after a specified record in a point file.

GENERAL FORM

CALL DKINS (RECID, AREA, L1, L2)

		•
where	RECID	is the logical record
		identifier after which
		the new record is to
		be entered. (INPUT)
		Is the record id of
	1251	the new record; (OUTPUT)
	AREA	is the first word of
		the information to be
		placed into the new
		record. (INPUT)
	L1	is the first character
	C1	
		in the record to receive
•		the first character in
		AREA. (INPUT)
	L2	is the last character
		in the record to receive
		information from AREA.
		(INPUT)

FIELD OPTIONS

FIELD	OPTIONS	DEFAULT
1. RECID	a. record id after	No Default

FIELD OPTIONS

FIEL	<u>.D</u>	OPTIONS	DEFAULT
		which the new re- cord will be placed b. zero to indicate placement as the first record in the file. c. computer output. record id of new data record.	1.
2.	AREA .	the location in core from which the writing will start.	!
3.	L1	 a. is the first character in the record to receive the first character in AREA. b. zero to indicate start with the first chracter. 	No Default
4.	L2	is the last charac- ter in the record to receive informa- tion from AREA.	

SPECIAL NOTES

The programmer must set NDS in common. If the programmer is writing a new record, FILENAME must be set in common. The programmer must define RECID, AREA, L1 and L2 before calling this subroutine.

EXECUTION PROCEDURES

A new record will be created after record RECID. If RECID is zero, a new record will be added to the beginning of the file. The system will place the first character in AREA in the L1 character location in the new record. Characters will be copied from AREA until the L2 character has been written into the new file. The record id for the new record will be returned in RECID.

OBTAIN LOGICAL RECORD LENGTH DKLEN

PURPOSE

To obtain the logical record length of an existing record.

GENERAL FORM

CALL DKLEN (RECID)

where RECID

is an existing blogical record id. (INPUT)

FIELD OPTIONS - None

SPECIAL NOTES

The programmer must set NDS in COMMON before calling this routine.

EXECUTION PROCEDURES

The system will obtain the length of record RECID in data set NDS and return it in variable LRLEN in common. The system sits PREID and FOLID also.

To make a file unavailable for writing.

GENERAL FORM

CALL DKLOS

FIELD OPTIONS - None

SPECIAL NOTES

The programmer must set NDS and FILENAME in common before calling this routine.

EXECUTION PROCEDURES

All records in core for FILENAME are marked for writing and will be placed onto the disk NDS. The file will be marked as closed and no one will be allowed to write into the file until another DHOPN is issued.

To find the data set number when the data set name is known.

GENERAL FORM

CALL (DSNAM, DSNO)

where

DSNAM

DSNO

is the name of .

the data set. (INPUT)

is the number of

the data set. (OUTPUT)

FIELD OPTIONS - None

SPECIAL NOTES - None

EXECUTION PROCEDURES

The system searches the data set name list and returns the data set number DSNO. If the dataset is not found, DSNO will be returned as zero.

DATA HANDLER INITIALIZATION DKNIT

PURPOSE

To initialize the data area in core for the data-handler and to establish the dataset name - logical reference number association.

GENERAL FORM

CALL DKNIT

FIELD OPTIONS - None

SPECIAL NOTES

This routine should be called once at the beginning of the application system.

EXECUTION PROCEDURES

The system will establish all in core data elements prior to the first use of any other data-handler subroutine.

To open a new or existing file.

GENERAL FORM

CALL DKOPN (IOP)

where IOP

is the storage option code for a new file only. (INPUT)

FIELD OPTIONS

FIELD	OPTIONS	DEFAULT
1. IOP	a. new files only 0 - neither point nor pack 1 - point only 2 - pack only 3 - point and pack b. not used for existing files	No Default

SPECIAL NOTES

The programmer must set NDS and FILENAME in COMMON. IOP is only applied if the file is new.

EXECUTION PROCEDURES

The file FILENAME is opened for writing in data set NDS using the original storage option given by the original DKOPN. If the file does not exist, the file is created.

<u>PURPOSE</u>

To write information to a record.

GENERAL FORM

CALL DKPUT (RECID, AREA, L1, L2)

where	RECID	is a record id in a data set. (INPUT) (OUTPUT)
	AREA	is the information to be written. (INPUT)
	L1	is the first character in the record to receive the
	L2	first character in AREA. (INPUT) is the last character in the record to receive data from AREA. (INPUT)

FIELD OPTIONS

FIELD	OPTIONS	DEFAULT	
1. RECID	a new reco end of a f (INPUT) b. a valid re	ile. cord id ting record. of the	

FIELD OPTIONS

FIELD		OPTIONS	DEFAULT
		(OUTPUT)	
2.	AREA	a. one computer word, variable name b. the first computer word to be used to hold the data. An element of a dimensioned array.	No Default
3.	L1 .	first character in the record to receive information.	No Default
4.	L2	 a. positive - last character in the record to receive i mation from AREA. b. negative - truncati of the logical recowill occur. The la character will be t absolute valve of L 	nfor- on rd st he

SPECIAL NOTES

The programmer must set NDS and FILENAME in COMMON.

EXECUTION PROCEDURE

If RECID is initially zero, a new record at the end of the file will be generated. The first character in AREA will be placed in the L1 character position in the record. Copying will continue until a character has

been written into the absolute value of L2. If L2 is negative, the existing record will be reduced to L2 characters.

To change the name of an existing file.

GENERAL FORM

CALL DKRNM (OCDNM, NEWNM)

where

OLDNM

is the current file name. (INPUT)

NEWNM

is the new file name. (INPUT)

FIELD OPTIONS - None

SPECIAL NOTES

The programmer must set NDS in COMMON before applying this routine.

EXECUTION PROCEDURE

The old name will be changed to the new name.

To obtain the current information related to a specific document.

GENERAL FORM

CALL DKSET (LBLK, NBU, NTALL)

where LBLK is the block length

in standard units. (OUTPUT)

NBU is the number of blocks

used. (OUTPUT)

NTALL is the total number of blocks.

used. (OUTPUT)

FIELD OPTIONS - None

SPECIAL NOTES

The programmer must set NDS before calling this routine.

EXECUTION PROCEDURE

The system will set all three parameters and return.

To clear all batters and write out all changes to the disks.

GENERAL FORM

CALL DKXIT

FIELD OPTIONS - None

SPECIAL NOTES

This routine should be called once before exiting from the application system.

EXECUTION PROCEDURES

The system will logically end all processing of data. All batters will be cleared and data change written to disk.

To obtain access to existing files.

GENERAL FORM

CALL LOCK (IND)
CALL UNLOK

where

IND

indicates control requested

FIELD OPTIONS

FIELD	OPTION	DEFAULT
1. IND	a. 0 - unconditional b - 1 - conditional	No Default

SPECIAL NOTES

During the actual execution of the data-handler in an interactive or multiprogramming environment, several users may be executing different copies of the code at the same time. Some of the users may wish to access (to read and/or write) the same files at the same time. The "multi-user" feature will permit concurrently executing programs to share the use of files. Shared use of these resources must be strictly controlled in order to ensure that one program does not interfere with the correct execution of others. This control takes the form of synchronizing use of these resources on the part of the programs invovled.

When a program requires use of a resource, it must request control of that resource from the operating system. Control of a resource can be either exclusive or shared. Exclusive control of a resource guarantees that no other program will be granted access to that resource (write access). Shared control guarantees that no other program will be granted exclusive control of that resource, but other programs will be granted shared control of that resource (read only access).

A request for either kind of control of a resource can be either conditional or unconditional. For conditional request, control is granted only if the resource is immediately available. The requesting program is informed as to whether or not control was granted. For unconditional requests, control is granted as soon as the resource becomes available. The requesting program may have to wait for an indefinite amount of time. This eventuality is entirely transparent to the program itself.

When a program has finished using a resource, it must relinquish control of that resource so that it becomes available for use by other programs.

REQUESTING AND RELINQUISHING CONTROL OF RESOURCES IS THE PROGRAMMER'S RESPONSIBILITY. Two subroutines, LOCK and UNLOK, have been made available for this purpose.

Required Programming Before Calling LOCK or UNLOK

INTEGER RSRCS COMMON /LOCKC/ RSRCS (7,40), NRSRCS

Before calling LOCK or UNLOK, you must describe the resources to be requested or relinquished via the array RSRCS. Each column corresponds to one resource. The number of resources described in RSRCS must be defined in the variable NRSRCS.

Before a call to LOCK, RSRCS(2,I), RSRCS(3,I), RSRCS(4,I), and RSRCS(5,I) must contain the name of the I'th resource requested. RSRCS(1,I) must contain the number of the data set on which this resource resides. RSRCS(7,I) should = 1 if exclusive control of this resource is requested, 0 if shared control is requested. NRSRCS should contain the number of resources, RSRCS(6,I) should not be used within a processing program.

LOCK, as its name implies, is to be used to request control of resources.

LOCK has a single argument, the integer variable IND, which should = 1 if the request is conditional, 0 if the request is unconditional. If IND = 1, LOCK will return the status of the request via the argument IND. IND = 1 if control of all requested resources was granted, 0 if control over none of the resources was granted due to the non-availability of one or more of the resources. If the request was unconditional, then control of all requested resources was granted (IND was not modified and still = 0). Note that LOCK (or UNLOK) never modifies LOCKC.

UNLOK is to be used to relinquish control of resources. UNLOK has no arguments, since control of resources is always relinquished unconditionally.

Application Rules

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The following rules must be strictly observed:

- (1) Control of a resource should not be requested until it is needed:
- (2) A resource must not be used until after its control has been requested and granted;
- (3) a. The use of DKOPN, DKPUT, DKINS, DKLOS, DKSET, DKRNM, DKXIT, or DKCLR in connection with a resource requires EXCLUSIVE CONTROL of that resource;
 - b. The use of DKFIL, DKLEN, or DKGET in connection with a resource requires shared or exclusive control of that resource:
 - c. The use of DKNDS, DKNIT, or DKDMP is unrestricted;
- (4) Control of a resource must not be relinquished until it is consistent with respect to both itself and other resources;
- (5) Control of a resource should be relinquished as soon as it is no longer needed.
- (6) A resource must not be used after its control has been relinquished;
- (7) The logic of execution should be such that: UNLOK is not

called without a previous call to LOCK; two calls to LOCK/UNLOK are never made without an intervening call to UNLOK/LOCK; LOCK is not called without a subsequent call to UNLOK.

- (8) a. The same resource must not be referenced by two different columns in RSRCS;
 - b. The maximum number of resources to be locked is 40;
- (9) No assumptions may be made regarding the contents of a resource at the time its control is granted.

Failure to observe one or more of the above rules will not necessarily result in the occurrence of a perceptible error condition (such as an ABEND), but resource integrity may nevertheless be seriously impaired. It is most important that great care in coding be exercised.

4 SUBROUTINE FUNCTIONS

Appendix B briefly describes the functions performed by each subprogram. This list is ordered by subroutine name.

Subroutines are divided into two general types: (1) application subroutines, and (2) function subroutines. The user's apply application programs application subroutines to request the data handler to perform work. The application and function subroutines use the function subroutines to actually perform work within the data handler.

5 ERROR MESSAGES

The following data-handler errors are defined below:

DATA-HANDLER PACKAGE ERROR NUMBERS

PAGE 1

ERROR NO.	ROUTINE	CONDITION
1	DKNIT -	SYST COMMON NOT INITIALIZED
2	DKNIT -	NOT ENOUGH SPACE FOR 1 BUFFER
3	DKGO -	NBSGN + NABLK .NE. UBU
4	DKODS -	OLD FILE NOT A DH FILE
2 3 4 5 6	DKODS -	CAN'T PROCESS OLD FILE - BLOCK SIZE TOO BIG
	DKODS -	ERROR READING OLD FILE HEADER
- 7	DKODS -	ERROR OPENING OLD FILE
8 9	DKODS -	ERROR OPENING NEW FILE
9	DKARG -	NDS INVALID
10	DKCDS -	ERROR CLOSING FILE & SAVING
11	DKCDS -	ERROR CLOSING FILE & DELETING
12	DKRBL -	READ INTO AN ACTIVE BUFFER - INTERNAL ERROR IN DH OR
		POLLUTED BUFFERS
13	DKASG -	CHARS IN FILE NAME MORE THAN ALLOWED
14	DKWAT -	INTERNAL ERROR - POLLUTED BT EMPTY - WAIT BIT IS ON
15	OKRBL -	I/O ERROR DURING READ
16	DKWBL -	WRITE FROM ACTIVE BUFFER
17	DKWBL -	I/O ERROR DURING WRITE
18	DKODS -	CAN'T PROCESS OLD FILE - NBSGNNBBLK INCONSISTENT

DATA-HANDLER PACKAGE ERROR NUMBERS (CONT'D)

ERROR NO.	ROUTINE	CONDITION
19 20 21	DKGET - GET WO	WITH THOSE ASSIGNED VIA DKNIT LENAME DOES NOT EXIST RK THAN MAYGP CHARS UED BIT NOT SET EXPECT TO READ MORE SEGMENTS-
22	DKNSG - CONTIN	LOGICAL INCONSISTENCY IN FILE UED BIT SET BUT NEXT LRID = OLOGICAL INCONSISTENCY IN FILE
23 24 25 26 27	DKSTO - PP HAS DKFND/DKGL2 - A DKGET - ILLEGAL DKRNM - INVALI	MENT NOT IN PP (NON-ZERO SEG COUNT IN HEADER) ZERO SEGS; ILLEGAL RECID ACCESS OF A DELETED LR SEGMENT L RECID (=0) D OLD FILENAME
28 29 30	DKPUT - FILE N	LRLEN+1 ON UPDATE (EXTEND) OT OPEN
31 32 33	DKPUT - INVALI	PUT MORE THAN MAXGP CHARS D FILENAME
34 35 36 37	DKDEL - LRID#O DKDEL - INVALIS).LT.L1 ON UPDATE FILE DOES NOT EXIST (RECORD DELETE) D FILENAME
38 39 40 41	DKDEL - PROTEC DKDEL - SHARED DKOPN - OPTION DKOPN - ILLEGA	DS OUT OF RANGE
42 43 44	DKOPN - OFT FUI	LL L NO EXTEND YET
45 46 47	DKOPN - SHAPED DKFIL - ILLEGAI DKGLS - ILLEGAI	DS L FILENAME L FILENAME
48 58 68	DKFIL - CAN'T (UNUSE - D	LENAME ALREADY EXISTS GET TO D.HAND TASK
69 70 71	UNUSE - D DKGET - L1.LT.	0 0
72 73 74 75 76	DKXDS - ALL BLOCK/UNLOK - NI	LKLEN P IS FULL. CANNOT BE EXTENDED. OCKS IN DATASET USED. RSRCS NOT POSITIVE. M IN DATASET TABLE.

6 PROGRAM CONVERSION

The data-handler has been written to minimize the time required to convert the system from one computer to another. All conversion instructions are given as comment statements within the code of the following subroutines:

DKBEX DKBIN DKCDS DKDDF DKERR DKG01 DKJCL DKMVC **DKNCS** DKNDS DKNIT DKODS DKPAK DKRBL DKTRC DKUPK DKWAT DKWBL DKXDS SNAP TIMDAT

The programmer performing this conversion should follow the instructions given in Appendix C. Several subroutines have been written in IBM BAL to improve executive speeds. These routines must be rewritten if the data-handler is being converted for use on computer systems other than the IBM 360/370 series. The following is a list of the applicable routines:

DKBEX DKBIN DKDDF DKJCL DKMYC DKNCS DKPAK DKUPK TIMDAT After each program has been rewritten in the new assembly language, the code should be added to the documentation tapes immediately after the previous version. The previous code should not be removed from the documentation tapes. This documentation system will allow the data-handler to be loaded in a relatively short time period onto a new computer system that is using a previously used assembler with minimal resources.

7 DOCUMENTATION TAPES

The documentation tapes provide a complete historic record of all computer programs ever written for the data-handler. All tapes are unlabeled, recorded at a density of 800 bits per inch (BPI), and blocked 800 characters per physical record. Each logical record is 80 characters in length.

The programs are listed in alphabetical order. Assembly programs are given in chronological order. The computer system and operating system version should be recorded in each assembly program in the first few comment cards.

All assembly programs are saved to allow conversion to a new computer assembly language with minimal resource expenditures. Each program is written as a separate tape file. The first tape file contains an index of all subsequent files. The format of this index is shown below:

DATA-HANDLER TAPE FILE INDEX

FILE	PROGRAM	SOURCE	DATE	DATE	LAST
NO.	NAME	LANGUAGE	WRITTEN	MODIFIED	
310 Right Justified	1320 Left Justified	2340 Left Justified	4350 DD/MM/YY	5360 DD/MM/YY	Card Columns Data Form

8 TEST ROUTINE

The last files on the documentation tapes contain a test routine that checks each application subroutine field. The correct results of this test routine are given on the tape for checking purposes.

9 CHANGES TO DATA-HANDLER PROGRAMS

As new functions are implemented into the data-handler, the following steps should be rigidly followed. Execution of these steps will insure that the changes are performed correctly and documented completely.

STEP ONE - Write a FORTRAN program, well documented with comment statements. Follow the standard documentation procedures below. If an existing program is being changed, change the FORTRAN program before changing the assembly programs.

Program Name: written as in the first line of the function or subroutine, with arguments.

<u>Function</u>: brief description of the function of the program.

Author: author's name, date of first writing.

Modifications: author, date and brief description of the reason for the change(s).

Language: programming language.

Calling Sequence: description of arguments and function results.

Routines Called: list of subroutines and functions called by the program.

Tasks or Modules: tasks or modules containing this routine.

Variables: a description of all variables used in the program.

COMMON variables will generally be described in a "main" program, referenced in this program. The structure of arrays and meaning of variables should be fully described.

Program Logic: A detailed description of the algorithm(s) used and the flow of the program. The comments generating this portion of the documentation will normally be scattered throughout the program in complete English sentences.

STEP TWO - Change the last date modified on the data-handler tape file index.

STEP THREE - Change the FORTRAN compile listing on the documentation tapes.

STEP FOUR - (If required) - Write the assembly programs.

 $\underline{\mathsf{STEP}}$ FIVE - (If required) - Change the last date modified on the data-handler tape file index.

STEP SIX - (If required) - Change the assembler listing on the documentation tapes.

STEP SEVEN - Change the test routines to completely test the new feature.

STEP EIGHT - Execute the entire test routine package and replace the results on the documentation tapes.

STEP NINE - Revise this report as required. Document all revisions listed for this document at the beginning of the report.

 $\underline{\mathsf{STEP}}$ TEN - Revise all application programs as required by the change in the data-handler.

APPENDIX A

DATA HANDLER

DUMP EXAMPLE

		•				
AEG. 1	00000000	00000000	0000000	00000004	00000000	00004F68
REG. 0	FFFFFDF4	FFFFFDF4	FFFFFF	00000000	00000000	00000144
REG. 14 HEG. 15 REG. 0	001119880	00215180	00208080	002126F0	00000000	000239F0 001406F8
REG. 14	42215244	422080E4	45212674	420DCOBE	421407F0	000239F0
TRACEBACK KOUTINE CALLED FROM ISH	0003	9400	900	0003	8000	
EBACK ROUTING	JRACE	DATHC	PAENR	LAPUI	16311	EDISPC
TRA						

AAAAERROK 30 IN DATA HANDLER PACKAGE.

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000000000000000000000000000000000000000	00000000	0000000	000000	00000000	10000000	90000000	00000000	00000001	00000014	00000000	00007FFF	0000000
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000000000000000000000000000000000000000	6FFFFFFF	00000000	00000000	80000000	04000010	00000000	000000E	00000000	00000000	00000000	00000000	30000000
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C====FILNM=====> FULID PREID LRLF4 FSTAT ====FILNM===== SYSTEM TABLE 00000000 U0H000000 U 1 40404040E2ESEZESCSD440E3

IPOOL IPLEH TOHER IDHER HUPBU LFIW LITHU MULM NBSGN NBBLK 201 9R00 0 0 30 1 32 32 256 8 24

***DATA SET TABLE IDSST **DA 1ELDS 1LDSN 200 12 12 2 ENTRY MU. 1 KDSLU <-----KDSFU----- KDSPR NUSPH KDSPL KDSSI KDSFD KDSBI KDSMN 00050001 AISYSTEM 256 1 0 254 20 21 105 1

OHOUGOUGO ENTRY MU. 2
ADSLU <----KUSFUS----- KDSPK AUSPU KUSPL KUSFU KUSFO KUSFS KDSBI KDSWN
0004000H BISYSIFF

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APPENDIX B: SUBROUTINE FUNCTION B.1 APPLICATION PROGRAMS

NUM- BER	SUBROU- TINE NAME	FUNCTION
1.	DKDEL	ROUTINE TO DELETE FILE FILNM (IF RECID=0) OR LR DECID IF NOT O FILE DOES NOT HAVE TO BE OPEN BUT MUST EXIST FOR RECORD DELETE FILE DELETE OF NON-EXISTENT FILE IS NOT AN ERROR
2.	DKFIL	ROUTINE SET FSTAT (COMMON) TO INDICATE STATUS OF FILMM (COMMON)
3.	DKGET	GET LOGICAL RECORD RECID IN DATA SET NDS (COMMON)
4.	DKINS	INSERT LOGICAL RECORD INTO DATA SET NDS (COMMON) IN FILE FILMM (COMMON) AFTER LR RECID.
5.	DKLEN	ROUTINE TO GET LENGTH OF LR RECID IN DATA SET NDS (COMMON)
6.	DKLOS	ROUTINE TO CLOSE FILE FILMM (COMMON) IN DATA SET NDS (COMMON)
7.	DKNDS	GIVEN DATA SET NAME, RETURN DATA SET NUMBER
8.	DKNIT	INITIALIZES ALL DATA HANDLER VARIABLES
9.	DKOPN	OPEN A FILE ROUTINE.
10.	DKPUT	PUT LOGICAL RECORD INTO DATA SET NDS (COMMON) IN FILE FILNM (COM)
11.	DKRNM	ROUTINE TO RENAME FILE
12.	DKSET	ROUTINE TO RETURN INFORMATION ABOUT DATA SET NDS
13.	DKXIT	ROUTINE TO LOOP THRU ALL DATA HANDLER DATA SETS (NDA) AND CLOSE EM DOWN

B.2 FUNCTION PROGRAMS

SUBROU-TINE

NAME FUNCTION

DKABP ADD TO BUFFER PRIORITY: CURRENTLY KEEPS AN ORDERED LIST OF BT ENTRIES BY MOST RECENT USE. N IS IGNORED.

DKALO ROUTINE TO TURN OFF THE PROTECT BIT FOR FILE FILMM IN DATA SET NDS. FILE NEED NOT BE OPEN

DKARG CHECK VALIDITY OF NDS AND OPEN DATA SET IF NOT ALREADY OPEN

DKASG ROUTINE TO ALLOW USER TO CONNECT A DATA SET NUMBER NSW TO A PARTICULAR LOGICAL UNIT NUMBERLUN AND FILE NAME FNAME.

DKBEX INTEGER FUNCTION TO EXTRACT NUMBR BITS FROM SOURCE WORD

DKBIN REAL FUNCTION TO INSERT THE LOW-ORDER NUMBR BITS FROM ISOR INTO RESLT AND RETURN IN REGISTER. RESIT IS NOT MODIFIED.

DKCBL ROUTINE TO COMPACT A BLOCK AT BUFFERIBTADD CALCULATE NUMBER OF SU'S TO MOVE FROM POOL SU OTH DISPL IFROM TO POOL SU OTH DISPL ITO BY SETTING NO. SU'S IN BLOCK FROM DST AND NO. SU'S FREE AT AND OF BLOCK (IN PR HEADER) MOVES SU'S ONE AT A TIME WITH DKPAK AND DKUPK AND UPDATES SU'S FREE AT END OF BLOCK. ASSUMES IFROM.GE.ITO EVENTUALLY MUST ALSO MODIFY FST

DKCDS OP SYS FILE INTERFACE ROUTINE TO CLOSE A DATA SET

DKDD1 GET ALL MACHINE CHARS IN MACH (1-13)

DKDD2 DATA SET DUMP ROUTINES

DKDD3 DATA SET DUMP ROUTINES

DKDD4 NOT WRITTEN

DKDDF ROUTINE FOR DYNAMIC DEFINE FILE STATEMENT IN OS FORTRAN ISSUES OPEN ON DDNAME FTNNFOO1

DKDDS DUMP DATA SET ROUTINE TO PRODUCE A FORMATTED DUMP OF DATA SET NDS

DKDLF ROUTINE TO DELETE ENTIRE FILMM (COMMON) IN DATA SET NDS (COMMON)

DKDL RROUTINE TO DELETE LR (NSGNBL) IN FILE FILMM DATA SET NDS

DKDMP DATA HANDLER

SUBROU-TINE

NAME FUNCTION

- DKDSG ROUTINE TO DELETE THE SEGMENT WHOSE LR HEADER IS IN H IN BUFFER IBTADD ASSUMING KBTPP SET BY DKSTQ (VIA DKFND) DOES NOT PRESENTLY COMPACT THE BLOCK OR MODIFY FST
- DKEBF ROUTINE TO EMPTY THE BUFFER AT IBTAD WHOSE HEADER IS IN D CHECKS FOR PREVIOUS I/O COMPLETION WRITES AND WAITS IF WRITTEN INTO AND ZERO ENTRY AND D
- DKERR CKECKS ERROR HANDLING SWITCH AND EITHER PRINTS AND ABORTS OR RETURNS WITHOUT SAYING ANYTHING
- DKFBF FILL BUFFER ROUTINE: BT ENTRY AT IBTAD IS IN D GET BLOCK NBL OF DATA SET NSW IN AND WAIT FOR READ COMPLETION
- DKFND ROUTINE TO LOCATE BLOCK NBL SEGMENT NSG OF DATA SET NSW.

 SEARCHES BT'S FIRST THEN READS INTO CORE RETURNS IBTAD=THE BT ENTRY
 DISPL WHERE PR IS DOES NOT SET MEMORY
- DKGBF GET A BUFFER ROUTINE: CHOOSES THE NEXT AVAILABLE BUFFER USING DKSBF ROUTINE SETS ITS PRIORITY HIGH WIHH DKABP AND EMPTIES IT WITH DKEBF. D AND BT ENTRY WILL BE ZEROED AND ALL IO FINISHED ON RETURN.
- DKGBL ROUTINE TO GET A BLOC5 FROM FREE SPACE TABLE OF DATA SET
 NDS WORD FROM BIT MAP STORED IN DST (IDSWD) NUMBER BST (IDSWN)
 BITS CORRESPOND TO BLOCKS NUMBER RT TO LEFT IN EACH WORD IF DST
 (IDSWN) = 0, NO WORD THERE. ALWAYS ASSUME WORD MUT BE REPLD
 USES INTERNAL ARRAY A TO STORE DST ENTRY
- DKGBM ROUTINE TO GET A BIT MAP WORD FOR NS WHICH HAS AT LEAST ONE BLOCK AVAILABLE (BIT=0) WITHIN THE BLOCKS INITIALIZED
- DKLG1 ROUTINE TO GET THE SEGMENT OF THE LR (NSWNSEGNBLK) CONTAINING NCS INTO CORE BUFFER AND SETS D=BT HEADER H=LR HEADER SETS QPPP FOR BUFFER CONTAINING DESIRED SEGMENT ALSO SETS PREID FOLID AND LRLEN IN COMMON FROM EITHER LAST BT MEMORY IF SEGMENT FOUND THERE OR FROM FIRST AND LAST BUFFER EXCEPT FOR THE LOOP TO GET LAST SEGMENT

化复数多位复数 医克尔克氏试验检尿 医动物

- DKGLS ROUTINE TO GET LAST SEGMENT OF LR (NSGNBL) IN NDS
- DKGO ROUTINE TO INITIALIZE THE DATA HANDLER BUILDS NECESSARY TABLES FROM ARRAY IN BLANK COMMON STARTING AT WORD ISTRT FOR ILEN WORS ALSO SETS CONSTANTS AND INITIALIZES VARIABLES
- DKGO1 SET MACHINE PARAMETERS IFR IBM 360/370 WITH OS

SUBROU-NUM TINE BER NAME FUNCTION

DKIDS ROUTINE TO INITIALIZE THE RECORDS IN A NEW DATA SET IN PARTICULAR TO SET UP BLOCK 1 WITH TWO SEGMENTS:

DKJCL

DKMFS ROUTINE TO SEARCH THE FST AND LOCATE THE BLOCK WITH THE MOST FREE SPACE WHICH IS RETURNED IN NBL.

DKMIC ROUTINE TO SEARCH IN CORE BUFFER TABLES FOR ONE WITH DATA SET NDS FILENAME FILMM AND RETURN THE ONE WITH THE MOST FREE SPACE.

DKMIR ROUTINE FOR BUFFER MEMORY ID RESET TO ENSURE THAT WHEN A NEW LOGICAL RECORD (RID) IS ADDED OR INSERTED THE PRECEEDING (NSGPNBLP) AND FOLLOWING (NSGFNBLF) RECORDS (IF THEY APPEAR IN ANY BT MEMORY) HAVE THEIR FOLID AND PREID ENTRIES RESPECTIVELY EQUAL RID.

DKMLR ROUTINE FOR MEMORY LENGTH RESET IN BT ENTRIES

DKMST SETS BUFFER MEMORY PARAMETERS INTO D

DKNCS ROUTINE TO COMPUTE THE NUMBER OF CHARACTERS IN A LR SEGMENTS

DKNSG ROUTINE TO GET NEXT SEGMENT OF THE LR WHOSE HEADER IS IN H INTO A BUFFER SETS BT ADDR IN IBTAD BT ENTRY IN D.

DKOUT BRINGS ACTIVITY ON THIS DATA SET TO AN ORDERLY HALT CLOSE ANY OPEN DH FILES IN THIS DATA SET AND WAIT FOR ANY BUFFER I/O ACTIVITY.

DKPA1 PUT ADD RECORD BUFFERING LOGIC - ACTION 1: LR WILL ENTIRELY FIT INTO IN-CORE BUFFER AT IBTADD

DKPA2 PUT ADD RECORD BUFFERING LOGIC - ACTION 2: LR PLACED IN REMAINDER OF THE IN-CORE BUFFER AT IBTADD AND THEN IN ONE OR MORE NEW BLOCKS. THIS LATTER IS DONE VIA DKPA4.

DKPA3 NOT WRITTEN YET

DKPA4 PUT ADD RECORD BUFFERING LOGIC - ACTION 4: LR MUST BE PLACED IN ONE OR MORE NEW BLOCKS

DKPA5 PROGRAM NOT WRITTEN YET

DKPAK ROUTINE TO PLACE ICT WORDS OF D (IN SU'S, NBU BITS) INTO POOL ARRAY

DKPAR ROUTINE TO DETERMINE THE PUT ADD RECORD BUFFERING RULE

SUBROU-

NUM TINE

BER NAME FUNCTION

DKPBL PUT BLOCK NBL BACK IN BIT MAP OF DATA SET NDS SO IT CAN BE REUSED

DKPBM ROUTINE TO REPLACE THE NTH SU (IN WORD) IN THE BIT MAP OF NDS

DKPMD(A) ROUTINE TO UPDATE THE PTR IN THE LRID WRITTEN PRIOR TO CURRENT ONE IF IN CHAIN MODE

DKPRO ROUTINE TO SET PROTECT BIT FOR FILE FILMM IN DATA SET NDS

DKPSG PUT LR SEGMENT ROUTINE USED TO ADD A NEW SEGMENT AT END OF BUF

DKPSH ROUTINE TO REPLACE THE LR SEGMENT HEADER IN H INTO BUFFER

DKRBL READS BLOCK NBL FROM DATA SET NSW INTO BT ENTRY IBTAD(IN D)

DKSBF SELECT BUFFER ROUTINE

DKSFD SEARCH FILE DIRECTORY ROUTINE IN DS NSW

DKSOF SEARCH OPEN FILE TABLE ROUTINE FOR FILMM (COMMON) IN DATA SET NDS (COMMON) RETURNS IOFAD:

DKSTQ ROUTINE TO LOOP THRU BT ENTRY IBTAD AND LOCATE SEGMENT NSG

DKTLR ROUTINE TO TRUNCATE CURRENT SEGMENT AT CHAR NCE COMPACT THE BLOCK AND DELETE ANY FOLLOWING SEGMENTS.

DKTRC ROUTINE TO PROVIDE A TRACEBACK AND OTHER INFO WHEN ERROR CALLED BY: DKERR IBM OS VERSION H LEVEL OBJECT TIME SYSTEM

DKUPK ROUTINE TO PLACE ICT SU'S OF POOL ARRAY P BEGINNING AT ISTRT+1 INTO D (ONE SU PER WORD)

DKWAT WAIT ROUTINE: WAITS FOR BIT IBIT TO CLEAR IF SET IN BT STATUS WORD. ALLOWS ASYNCHRONOUS I/O OPERATIONS WITH THE PROPER OP SYS INTERFACE SUPPORT (DKRBL AND DKWBL).

DKWBL WRITES BLOCK NBL FROM BT ENTRY IBTAB(IN D) INTO DATA SET NSW

DKWIN ROUTINE TO SET THE 'WRITTEN INTO' BIT IN BT ENTRY IN D
DOES NOT REPLACE D

DKXBM ROUTINE TO EXTEND BIT MAP BY ONE SEGMENT

DKXDS ROUTINE TO EXTEND DATA SET

SUBROU-

NUM TINE BER NAME

FUNCTION

DKXFD EXTEND FILE DIRECTORY OF DATA SET NDS (COMMON) BY ONE BLOCK

DKXLR ROUTINE TO EXTEND THE LR (NSGNBL) WHOSE LAST SEGMENT IS IN BT ENTRY IBTAD.D AND WHOSE LR HEADER IS IN HEDLR

XOR

DKODS OPEN DATA SET TO HOST OP SYSTEMS FILE CONTROL SERVICES

DKDMP DUMPS DATA HANDLER COMMON AND POOL THIS ROUTINE IS A FUNCTION OF MACHINE CHARACTERISTICS SUCH AS CHARS PER WORD AND CHARS IN A DATA SET NAME

GOT236 THIS ROUTINE IS ENTERED BY THE FORTRAN EXTENDED ERROR HANDLING ROUTINES

BITEX RIGHT JUSTIFY THE BITS TO BE RETURNED

BITIN THE RIGHTMOST NUMBR BITS OF SOURC ARE INSERTED INTO RESLT BEGINNING WITH BIT START OF RESLT. THE REMAINING BITS OF RESLT ARE NOT MODIFIED

ICOPY COPIES NUM CHARACTERS FROM SOURC (STARTING AT SFST) TO DEST (STARTING WITH DFST) ALL OTHER CHARACTERS OF DEST ARE UNCHANGED

PACK NUM CHARACTERS INTO THE STRING DEST STARTING WITH CHARACTER L2 OF DEST. THE CHARACTERS ARE TAKEN FROM THE LEAST SIGNIFICANT CHARACTER OF THE ELEMENTS OF THE ARRAY SOURC STARTING WITH STANDARD UNIT L1.

SNAP

SYSCM INITIALIZE SYSTEM COMMON TO THE STANDARD VALUES REQUIRED BY THE PARTICULAR MACHINE IMPLEMENTATION

DKMVC MOVE NC CHARACTERS TO A(TO)+ITOFF
DKBNF ROUTINE TO INSERT AN INTEGER FORTRAN WORD INTO A REAL FORTRAN WORD

DKBXF ROUTINE TO EXTRACT AN INTEGER FORTRAN WORD FROM A REAL FORTRAN WORD.

AND

DKCMP

SUBROU-NUM TINE BER NAME FUNCTION

OR

SHIFT

TIMDAT

4

